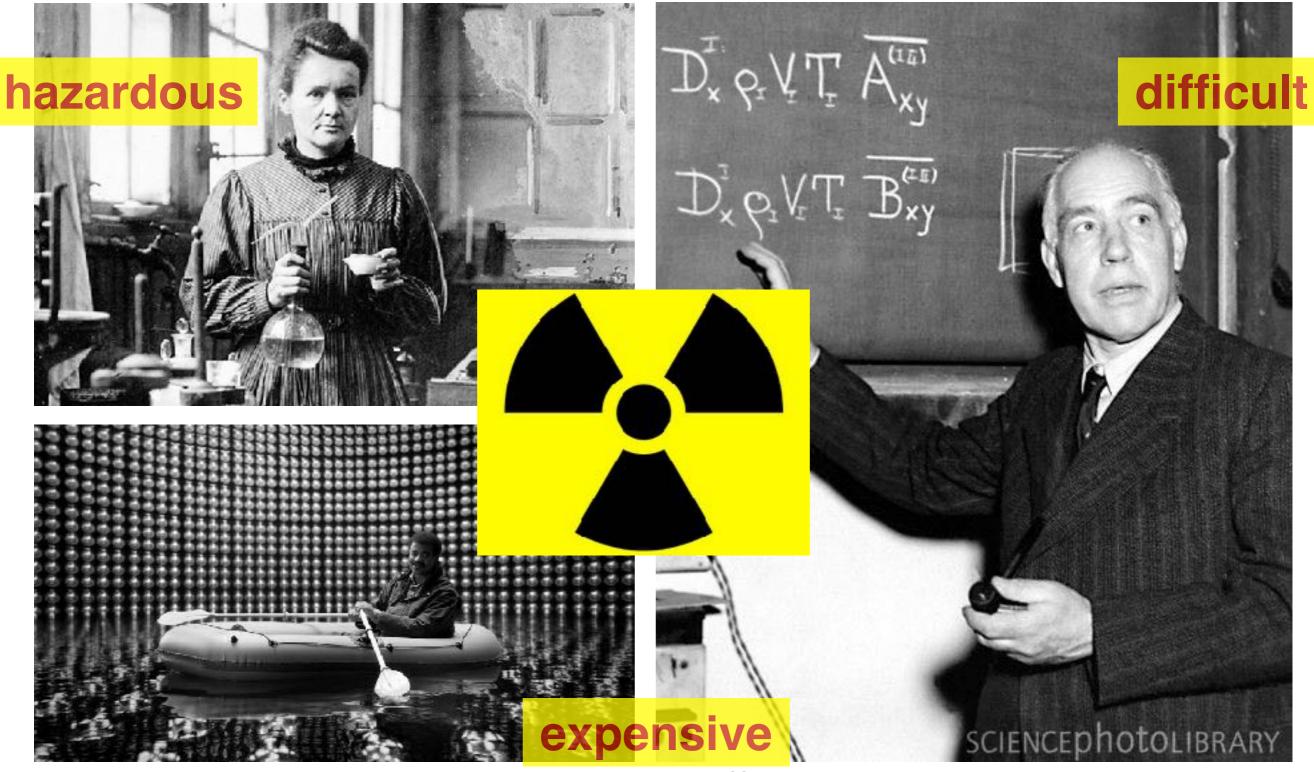
## Online Processing p.1 Introduction

Radoslaw Karabowicz Denis Klein FairRoot Group, GSI

#### Nuclear experiments



## Computing challenges

new HEP experiments producs huge amounts of data



traditional trigger systems are extended by heterogenous (CPU, GPU, FPGA, ...) online computing farms to reduce the data flow



the stored data are repeatedly being analyzed in the offline computing facilities (farms, grid, cloud) Radoslaw Karabowicz, GSI



#### Software levels

- Data transport
- Farm configuration
- Farm control & management
- Data reconstruction
- Physics analysis
- Parameter accessibility
- Parameter updates

#### Framework requierements

- Extend traditional single-process FairRoot to allow simulation and reconstruction of streaming data.
- High performance, flexibility and scalability are required!
- Separate running of tasks in **multiple processes** (that can be multi-threaded or use GPU) to increase stability (process crashes do not immediately affect other components and allow easier recovery).
- Communicate via **platform-independent messages** (simplify integration of different hardware and programming languages).

#### FairMQ

- Project started in 2012.
- People involved: Mohammad Al-Turany, Dennis Klein, Alexey Rybalchenko, Nicolas Winckler, Dennis Klein.
- Message Queue (MQ) transport messages between processes using queues.
- FairMQ data transport framework to organize processing tasks in topologies, consisting of independent processes (devices), communicating via asynchronous message queues over network or inter-process channels.

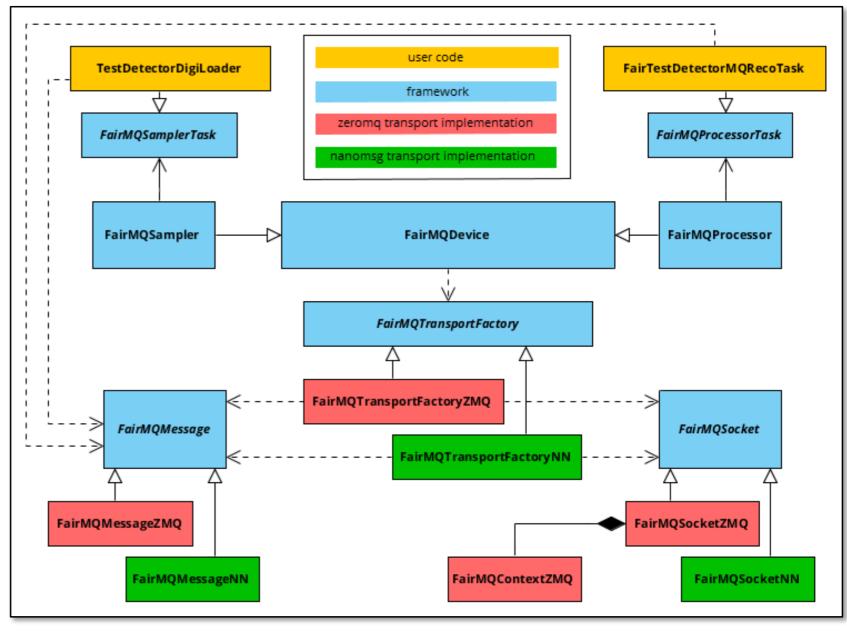
# FairMQ Transport Interface

FairMQ transport interface keeps the device and task code independent of the transport implementation (implementing Send/Receive/Polling methods, (Multipart) messages, patterns).

Currently two implementations:

- **ZeroMQ** (TCP, IPC, INPROC, PGM, TIPC, VMCI);
- nanomsg (TCP, IPC, INPROC, WebSockets).

Open for implementation with future technologies. Has to fit in the message passing model (or interface it).



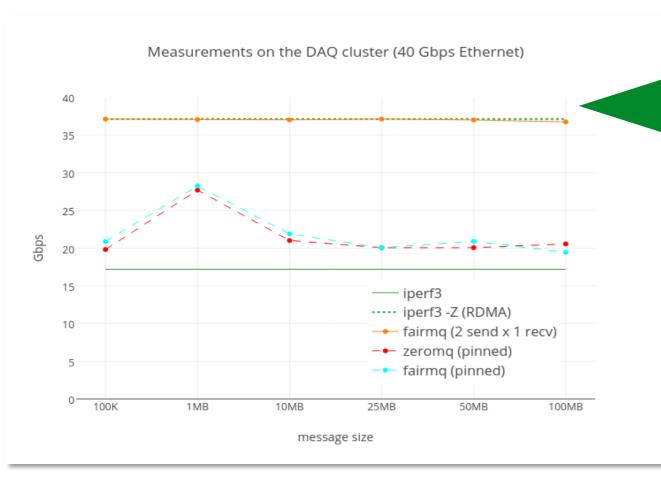
#### FairMQ

Transport based on:

ZeroMQ http://zeromq.org/ nanomsg http://nanomsg.org/ // production quality, large community, active development, extensive documentation, examples and patterns

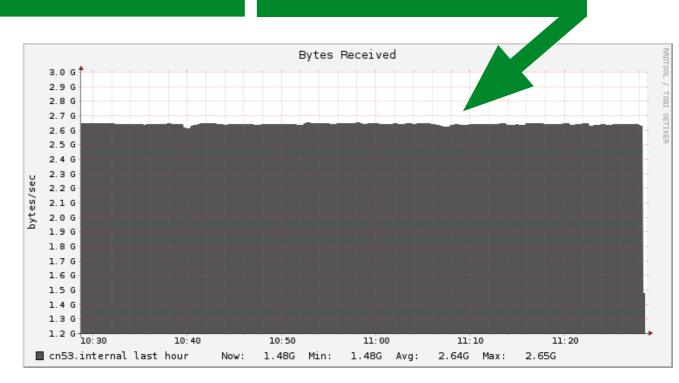
// still beta, several architectural improvements, cleaner internal API (eg. for adding new transport technologies)

Low overhead abstractions, staying close to theoretical performance limit of the protocol/network technology



2x Intel Xeon E5-2650 @ 2.60GHz (8 cores, 16 threads) (dual socket), 40 Gbps Ethernet, 3.10.x kernel

#### Sustained high performance on Ethernet, InfiniBand (IPoIB)



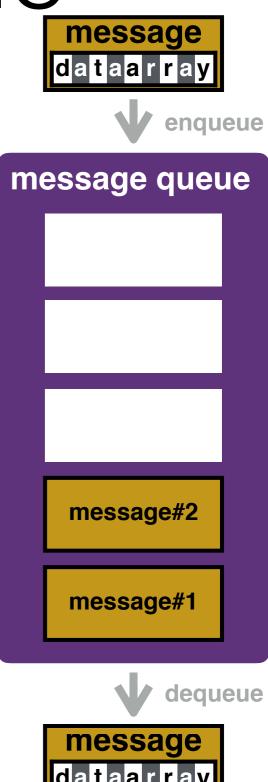
2 x Intel Xeon E5520 @ 2.27GHz (4 cores each, 8 threads) (dual socket), QDR InfiniBand (IPoIB), 3 Processes on 3.17.x kernel.

# FairMQ general

- Promotes separation of concerns to increase maintainability and reduce coupling: simple single-responsibility devices.
- Keeps the user and transport code separate, allowing them to evolve independently: abstract transport interface, user task separation.
- Simplifies and unifies handling of transport details: socket settings, polling, termination, access to communication patterns.
- Takes care of the device state (change, query), termination, command interface, device configuration, logging.

# Message Queue

- Message data to be transported between two processes, essentially it is an array of bytes (defined by the address to the first byte and the size of the array) and a header:
  - FairMQMessage simple message;
  - FarMQParts message composed of multiple parts.
- Queue array of messages waiting to be processed, arranged in a FIFO pipeline.



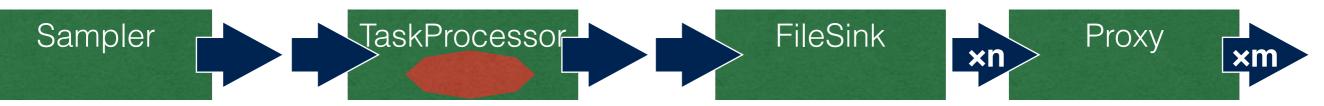
### FairMQ task

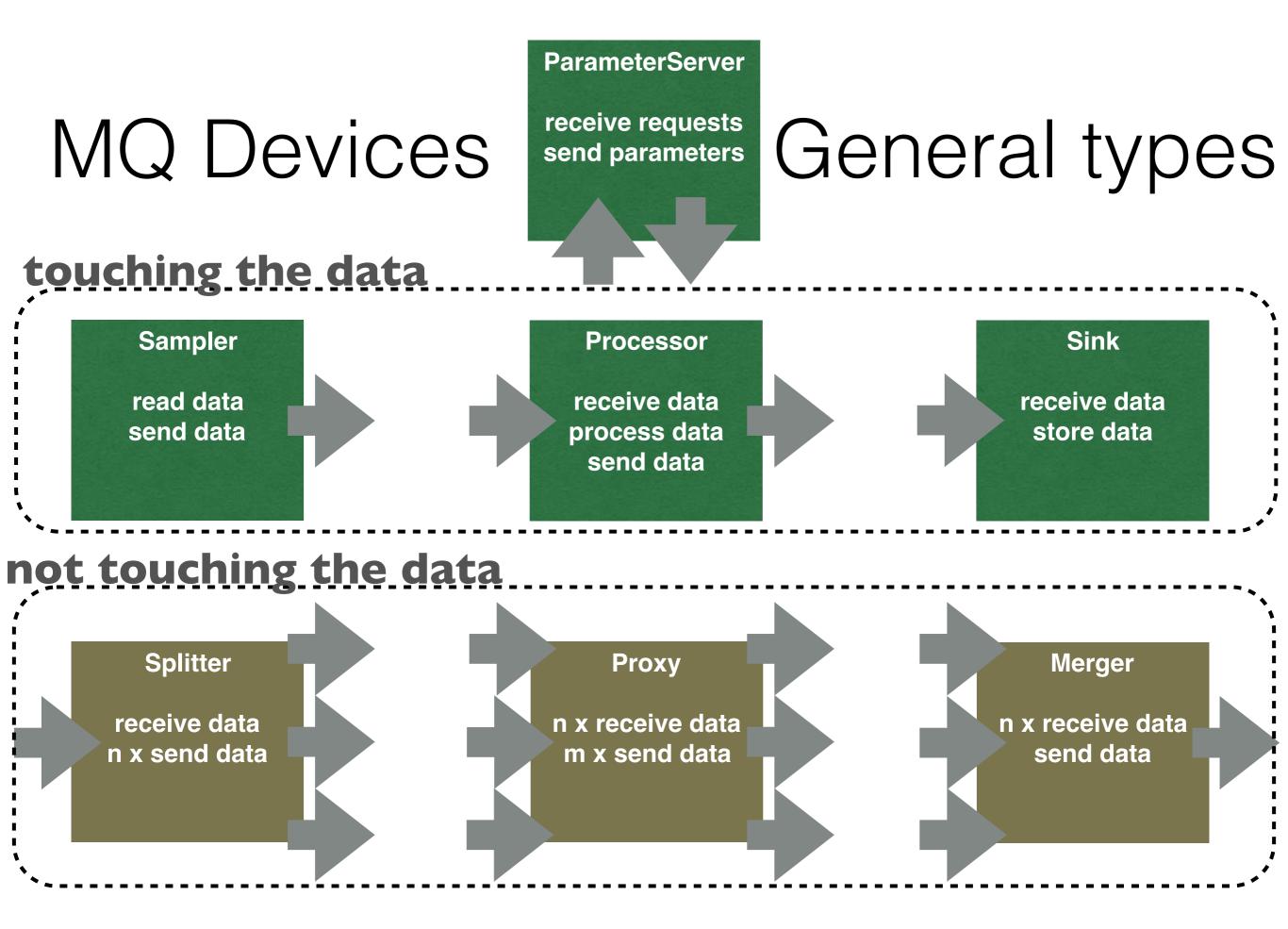
 Task contains an algorithm to address (hopefully) one specific physics problem (f.e. cluster finding, track fitting, PID assignment, etc.)

- It relates directly to FairTask.
- But, contrary to running under FairRoot, in FairMQ the task has no direct access to FairRunAna, FairRootManager, FairRuntimeDb

## FairMQ devices

- Simple single-responsibility computer programm able to receive and send messages over various channels.
- Messages may contain data, runtime parameters, configuration or communication data.
- It may encapsulate a task.
- There are also devices not connected to tasks, for example to provide IO access, to read and send parameters, to split or merge data streams.





## FairMQ channel

- Channel connects devices.
- There is always one device that binds a specific port on the host machine.
- Another devices connect to that port at a given address.
- The messaging pattern defines the way of message exchange.

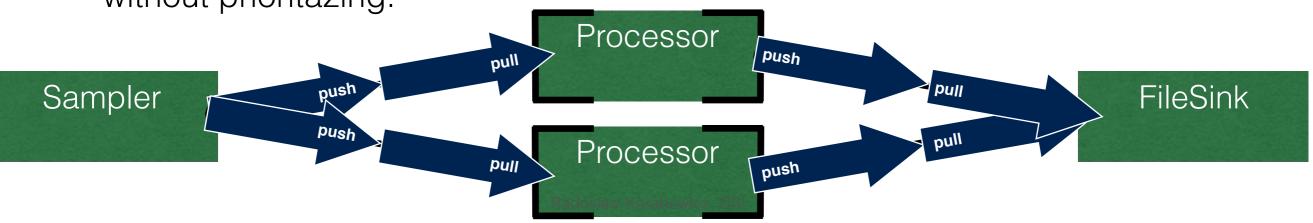


# FairMQ req-rep

- the common example of request-reply pattern is the comunication between a server and clients;
- the clients **request** a specific data or action;
- the server **replies** by sending the data or an answer;
- in this pattern the server has to bind, the clients connect.
  Processor
  Processor
  Processor
  Processor
  Processor

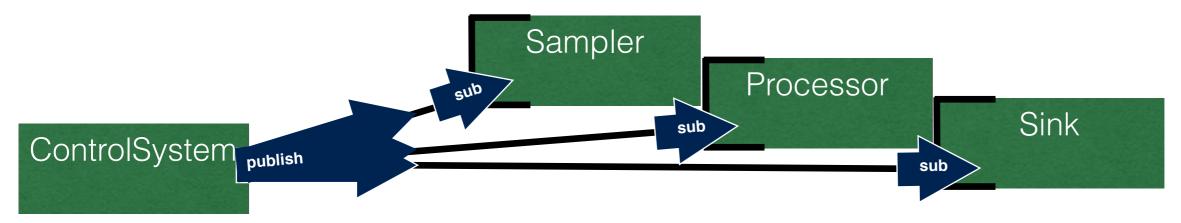
# FairMQ push-pull

- the push-pull pattern works like a pipeline: the messages flow always in one direction;
- one side **pushes** the messages into the queue;
- another side **pulls** the messages from the queue;
- this pattern works either 1to1, 1toN or Nto1;
- in case of 1toN, the messages will be distributed among the N receivers evenly;
- in the case of Nto1, the receiver will receive the messages from senders without prioritazing.



# FairMQ pub-sub

- the publish-subscribe pattern could be used for controling purposes;
- the subscriber(s) gets the messages after they are published by the publisher(s);
- it works 1to1, 1toN or Nto1;
- if no subscribers connected, the messages will be lost;
- otherwise each subscriber gets each message.



#### FairMQ PAIR

• bidirectional communication between two nodes.

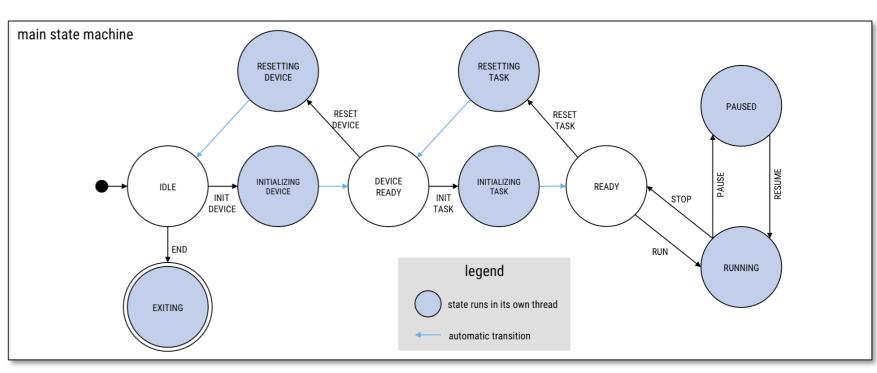


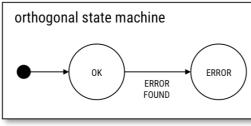
## FairMQ State Machine

A complex topology may require some degree of **synchronization** between different stages of the processing pipeline.

To allow this, the state of the devices can be queried and changed in response to an external control system. This is handled by the **FairMQStateMachine**.

The **device remains responsive** regardless of the condition of the user code, by running it in a separate thread.





Additionally, an orthogonal OK/ERROR state reflects the device condition. If error occurs, the device remains in the ERROR state to allow debuggin.

# FairMQ topology

- the different devices are used as "LEGO" blocks;
- the structure of all the blocks used for a given problem is called a **topology**;
- it lists the devices used, states the connections between them and describes the communication patterns;

# FairMQ topology

- defined in a JSON file;
- contains the list of devices (by the unique IDs);
- specifies connection types and communication patterns, addresses and ports, buffer sizes, verbosity level;
- example view of a topology with 1 sampler (bind), 1 sink (bind), 3 processors (connect) with push-pull communication pattern + a patameter server accessed via request-reply:

